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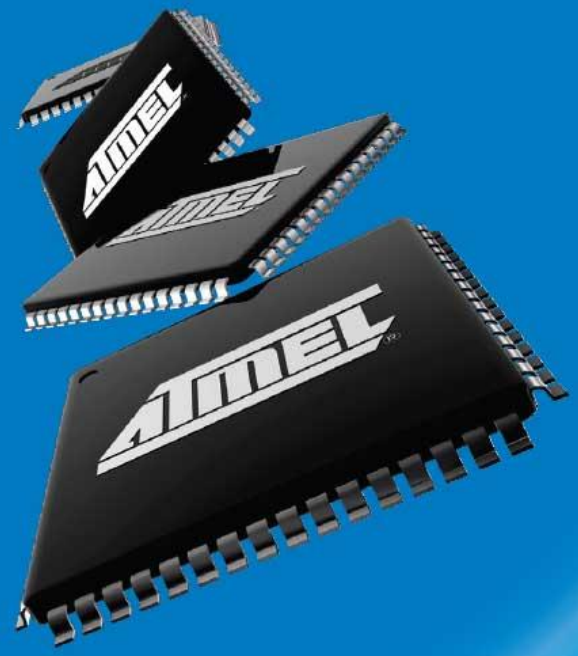


Addressing Modes

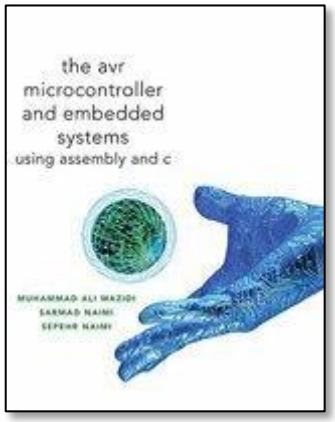
February 2009

Everywhere You Are®

Addressing Modes Part I – Working with AVR's Load-Store RISC Architecure

READING

The AVR Microcontroller and Embedded Systems using Assembly and C)
by Muhammad Ali Mazidi, Sarmad Naimi, and Sepehr Naimi



Chapter 2: AVR Architecture and Assembly Language Programming

Section 2.3: Using Instructions with the Data Memory

Chapter 6: AVR Advanced Assembly Language Programming

Section 6.2: Register and Direct Addressing Modes

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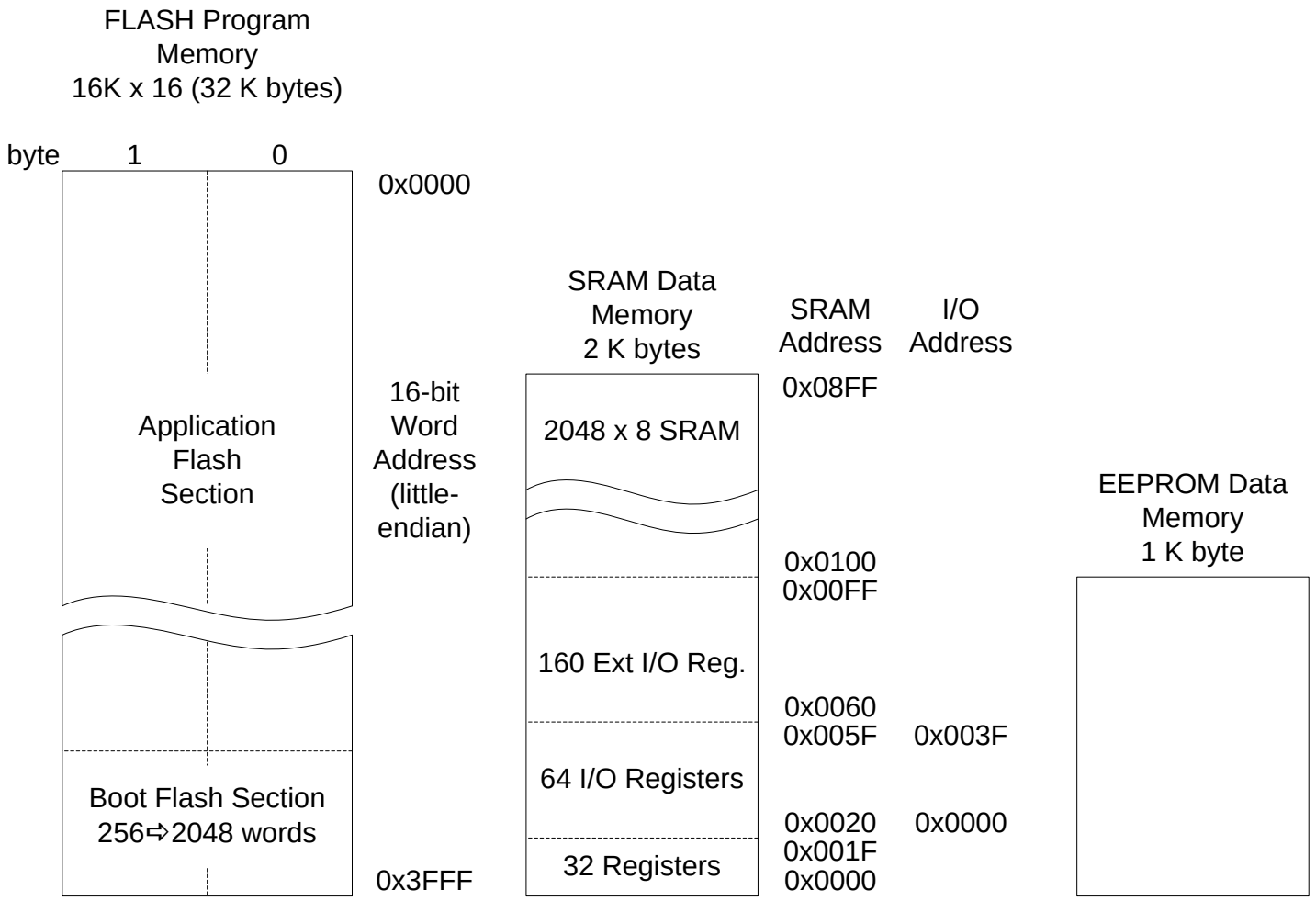
LOAD-STORE INSTRUCTIONS AND THE ATMEGA328P MEMORY MODEL

- When selecting an addressing mode you should ask yourself where is the operand (data) located within the memory model of the AVR processor and when do I know its address (assembly time or at run time)[1].



[1] http://www.atmel.com/dyn/resources/prod_documents/doc0856.pdf 8-bit AVR Instruction Set

LOAD-STORE INSTRUCTIONS AND ADDRESSING MODES

- When loading and storing data we have several ways to “address” the data.
- The AVR microcontroller supports addressing modes for access to the Program memory (Flash) and Data memory (SRAM, Register file, I/O Memory, and Extended I/O Memory).

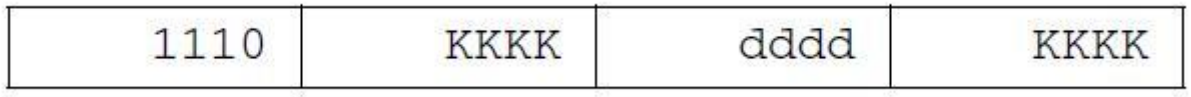
Load-store Instructions

	Address Space		
Addressing Mode	Flash Program	SRAM Data	I/O
Immediate	`ldi`		
Direct		`lds, sts`	`in, out`
Indirect	`lpm, spm (1)`	`ld, st (2)`	
Indirect with Displacement		`ldd, std (3)`	

IMMEDIATE

- Data is encoded with the instruction. Operand is therefore located in Flash Program Memory. This is why technically our memory model is a *Modified* Harvard.

```
ldi    r16, 0x23      // where ldi = 1110, Rd = 0000₂, and constant K = 00100011₂
```

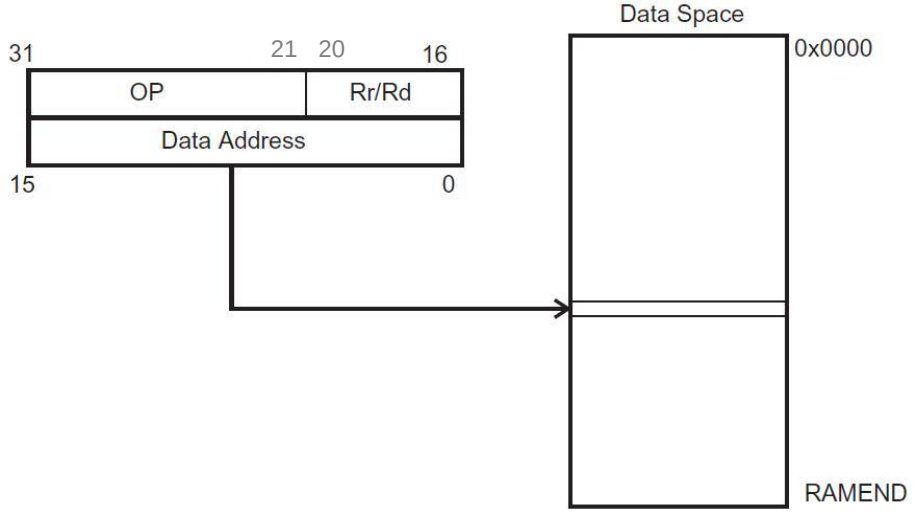


- Notice that only four bits (dddd) are set aside for defining destination register Rd. This limits us to 2^4 = 16 registers. The designers of the AVR processor chose registers 16 to 31 to be these registers (i.e., 16 ≤ Rd ≤ 31).

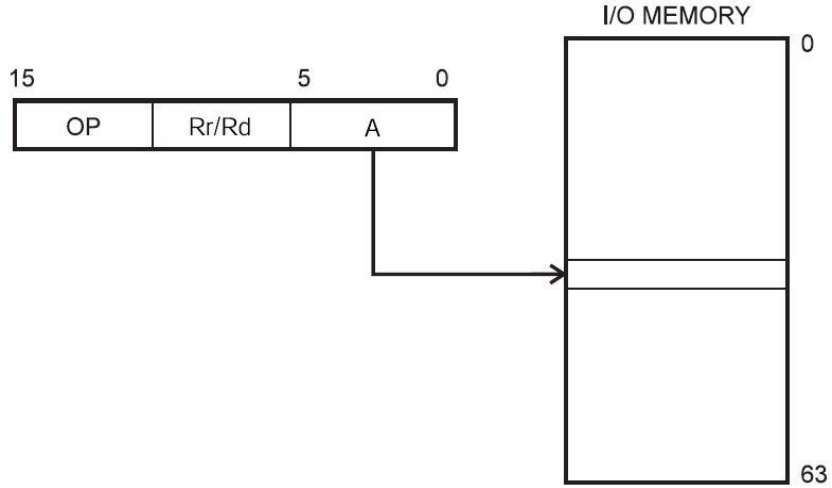
DIRECT

```
lds     r16, A
sts     A, r16
```

Within the AVR family there are two (2) possible lds/sts instructions. A specific family member will have only one lds/sts combination. The ATmega328P lds/sts instruction is illustrated here with the exception that 5 bits (not 4) encode Rr/Rd. This means all 32 registers are available to the lds/sts instruction.




```
in      r16, PINC
out     PORTD, r16
```



REGISTER-REGISTER INSTRUCTIONS

Data Transfer

- Register-register move byte (`mov`) or word (`movw`)

Arithmetic and Logic (ALU)

- Two's complement negate (`neg`), Arithmetic add (`add, adc, adiw`), subtract (`sub, subi, sbc, sbci`), and multiply (`mul, muls, mulsu, fmul, fmuls, fmulsu`)
- Logical not (`com`), and (`and, andi, cbr, tst`), or (`or, ori, sbr`), exclusive or (`eor`)
- Clear (`clr`), set (`ser`), increment (`inc`), decrement (`dec`)

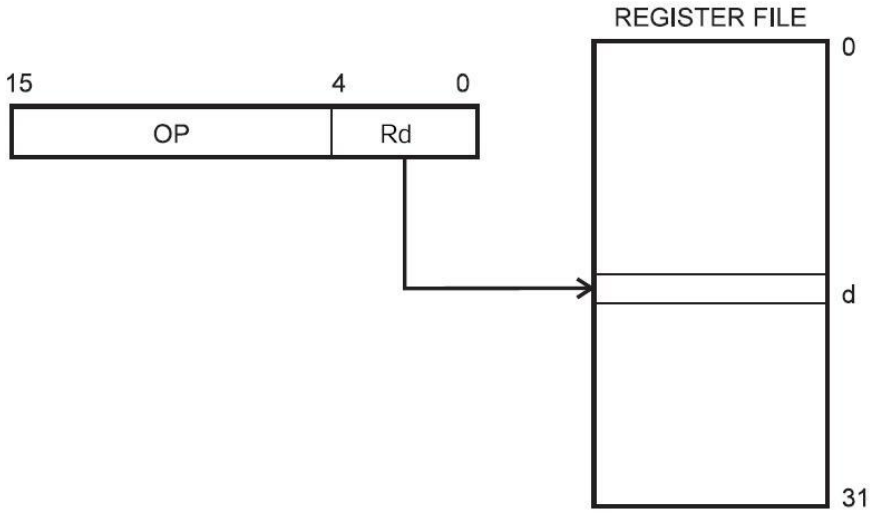
Bit and Bit-Test

- Register logical shift left (`lsl`) or right(`lsr`); arithmetic shift right (`asr`); and rotate left or right (`rol`, `ror`)
- Register swap nibble (`swap`)
- Register bit load (`bld`) or store (`bst`) from/to T flag in the Status Register SREG
- I/O Register Clear (`cbi`) or set (`sbi`) a bit
- Clear (`cl`*`Flag`*) or set (`se`*`Flag`*) a Flag bit in the Status Register SREG by name (I, T, H, S, V, N, Z, C) or bit (`bclr`, `bset`).

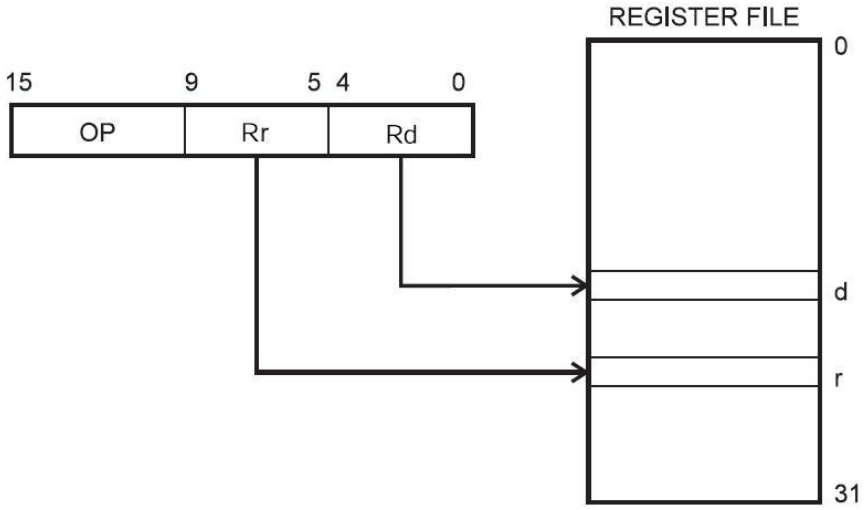
REGISTER DIRECT

In the following figures, OP means the operation code part of the instruction word. To simplify, not all figures show the exact location of the addressing bits. To generalize, the abstract terms RAMEND and FLASHEND have been used to represent the highest location in data and program space.

```
com     r16
```




```
add     r16, r17
```



LOAD-STORE PROGRAM EXAMPLE

Write an Assembly program to add two 8-bit numbers.

C = A + B

```
lds   r16, A      ; 1. Load variables
lds   r17, B
add   r16, r17    ; 2. Do something
sts   C, r16      ; 3. Store answer
```

- Identify the operation, source operand, destination operand in the first *Data Transfer* instruction.
- Identify the source/destination operand in the *Arithmetic and Logic* (ALU) instruction.
- What addressing mode is used by the source operand, in the first instruction?
- Show contents of Flash Program Memory (mnemonics)
- Show contents of SRAM Data Memory, assuming variables are stored in sequential memory locations starting at address 0100_{16}.
- Modify the program to leave register r16 unchanged by making a copy (use r15).

SPECIAL TOPIC – HARVARD VERSUS PRINCETON ARCHITECTURE

Princeton or Von Neumann Memory Model

Program and data share the same memory space. Processors used in all personal computers, like the Pentium, implement a von Neumann architecture.

Harvard Memory Model

Program and data memory are separated. The AVR processors among others including the Intel 8051 use this memory model. One advantage of the Harvard architecture for microcontrollers is that program memory can be wider than data memory. This allows the processor to implement more instructions while still working with 8-bit data. For the AVR processor program memory is 16-bits wide while data memory is only 8-bits.

You may have already noticed that when you single step your program in the simulator of AVR Studio it is incremented by 1 each time an instruction is executed. No surprise there right? Wrong. The program memory of the AVR processor can also be accessed at the byte level. In most cases this apparent paradox is transparent to the operation of your program with one important exception. When you want to access data stored in program memory, you will be working with byte addresses not words (16-bits). The assembler is not smart enough to know the difference and so when you ask for an address in program memory it returns its word address. To convert this word address into a byte address you need to multiply it by 2. Problematically we do this by using the shift left syntax of C++ to explicitly tell the assembler to multiply the word address by 2. Remember, when you shift left one place you are effectively multiplying by 2.

With this in mind, we would interpret the following AVR instruction as telling the AVR assembler to convert the word address of label beehives in program memory to a byte address and then to take the low order of the resulting value and put into the source operand of the instruction.

```
ldi     ZL,low(beeHives<<1)    // load word address of beeHives look-up
```

APPENDIX A – ATMEGA328P INSTRUCTION SET[2]

Mnemonics	Operands	Description	Operation	Flags	#Clocks
ARITHMETIC AND LOGIC INSTRUCTIONS					
ADD	Rd, Rr	Add two Registers	Rd ← Rd + Rr	Z,C,N,V,H	1
ADC	Rd, Rr	Add with Carry two Registers	Rd ← Rd + Rr + C	Z,C,N,V,H	1
ADIW	Rdl,K	Add Immediate to Word	Rdh:Rdl ← Rdh:Rdl + K	Z,C,N,V,S	2
SUB	Rd, Rr	Subtract two Registers	Rd ← Rd - Rr	Z,C,N,V,H	1
SUBI	Rd, K	Subtract Constant from Register	Rd ← Rd - K	Z,C,N,V,H	1
SBC	Rd, Rr	Subtract with Carry two Registers	Rd ← Rd - Rr - C	Z,C,N,V,H	1
SBCI	Rd, K	Subtract with Carry Constant from Reg.	Rd ← Rd - K - C	Z,C,N,V,H	1
SBIW	Rdl,K	Subtract Immediate from Word	Rdh:Rdl ← Rdh:Rdl - K	Z,C,N,V,S	2
AND	Rd, Rr	Logical AND Registers	Rd ← Rd • Rr	Z,N,V	1
ANDI	Rd, K	Logical AND Register and Constant	Rd ← Rd • K	Z,N,V	1
OR	Rd, Rr	Logical OR Registers	Rd ← Rd v Rr	Z,N,V	1
ORI	Rd, K	Logical OR Register and Constant	Rd ← Rd v K	Z,N,V	1
EOR	Rd, Rr	Exclusive OR Registers	Rd ← Rd ⊕ Rr	Z,N,V	1
COM	Rd	One's Complement	Rd ← 0xFF – Rd	Z,C,N,V	1
NEG	Rd	Two's Complement	Rd ← 0x00 – Rd	Z,C,N,V,H	1
SBR	Rd,K	Set Bit(s) in Register	Rd ← Rd v K	Z,N,V	1
CBR	Rd,K	Clear Bit(s) in Register	Rd ← Rd • (0xFF - K)	Z,N,V	1
INC	Rd	Increment	Rd ← Rd + 1	Z,N,V	1
DEC	Rd	Decrement	Rd ← Rd – 1	Z,N,V	1
TST	Rd	Test for Zero or Minus	Rd ← Rd • Rd	Z,N,V	1
CLR	Rd	Clear Register	Rd ← Rd ⊕ Rd	Z,N,V	1
SER	Rd	Set Register	Rd ← 0xFF	None	1
MUL	Rd, Rr	Multiply Unsigned	R1:R0 ← Rd x Rr	Z,C	2
MULS	Rd, Rr	Multiply Signed	R1:R0 ← Rd x Rr	Z,C	2
MULSU	Rd, Rr	Multiply Signed with Unsigned	R1:R0 ← Rd x Rr	Z,C	2
FMUL	Rd, Rr	Fractional Multiply Unsigned	R1:R0 ← (Rd x Rr) << 1	Z,C	2
FMULS	Rd, Rr	Fractional Multiply Signed	R1:R0 ← (Rd x Rr) << 1	Z,C	2
FMULSU	Rd, Rr	Fractional Multiply Signed with Unsigned	R1:R0 ← (Rd x Rr) << 1	Z,C	2
BRANCH INSTRUCTIONS					
RJMP	k	Relative Jump	PC ← PC + k + 1	None	2
IJMP		Indirect Jump to (Z)	PC ← Z	None	2
JMP[1]	k	Direct Jump	PC ← k	None	3
RCALL	k	Relative Subroutine Call	PC ← PC + k + 1	None	3
ICALL		Indirect Call to (Z)	PC ← Z	None	3
CALL[1]	k	Direct Subroutine Call	PC ← k	None	4
RET		Subroutine Return	PC ← STACK	None	4
RETI		Interrupt Return	PC ← STACK	I	4
CPSE	Rd,Rr	Compare, Skip if Equal	if (Rd = Rr) PC ← PC + 2 or 3	None	1/2/3
CP	Rd,Rr	Compare	Rd – Rr	Z, N,V,C,H	1
CPC	Rd,Rr	Compare with Carry	Rd – Rr – C	Z, N,V,C,H	1
CPI	Rd,K	Compare Register with Immediate	Rd – K	Z, N,V,C,H	1
SBRC	Rr, b	Skip if Bit in Register Cleared	if (Rr(b)=0) PC ← PC + 2 or 3	None	1/2/3
SBRS	Rr, b	Skip if Bit in Register is Set	if (Rr(b)=1) PC ← PC + 2 or 3	None	1/2/3
SBIC	P, b	Skip if Bit in I/O Register Cleared	if (P(b)=0) PC ← PC + 2 or 3	None	1/2/3
SBIS	P, b	Skip if Bit in I/O Register is Set	if (P(b)=1) PC ← PC + 2 or 3	None	1/2/3
BRBS	s, k	Branch if Status Flag Set	if (SREG(s) = 1) then PC←PC+k + 1	None	1/2
BRBC	s, k	Branch if Status Flag Cleared	if (SREG(s) = 0) then PC←PC+k + 1	None	1/2
BREQ	k	Branch if Equal	if (Z = 1) then PC ← PC + k + 1	None	1/2
BRNE	k	Branch if Not Equal	if (Z = 0) then PC ← PC + k + 1	None	1/2
BRCS	k	Branch if Carry Set	if (C = 1) then PC ← PC + k + 1	None	1/2
BRCC	k	Branch if Carry Cleared	if (C = 0) then PC ← PC + k + 1	None	1/2
BRSH	k	Branch if Same or Higher	if (C = 0) then PC ← PC + k + 1	None	1/2
BRLO	k	Branch if Lower	if (C = 1) then PC ← PC + k + 1	None	1/2
BRMI	k	Branch if Minus	if (N = 1) then PC ← PC + k + 1	None	1/2
BRPL	k	Branch if Plus	if (N = 0) then PC ← PC + k + 1	None	1/2
BRGE	k	Branch if Greater or Equal, Signed	if (N ⊕ V= 0) then PC ← PC + k + 1	None	1/2
BRLT	k	Branch if Less Than Zero, Signed	if (N ⊕ V= 1) then PC ← PC + k + 1	None	1/2
BRHS	k	Branch if Half Carry Flag Set	if (H = 1) then PC ← PC + k + 1	None	1/2
BRHC	k	Branch if Half Carry Flag Cleared	if (H = 0) then PC ← PC + k + 1	None	1/2
BRTS	k	Branch if T Flag Set	if (T = 1) then PC ← PC + k + 1	None	1/2
BRTC	k	Branch if T Flag Cleared	if (T = 0) then PC ← PC + k + 1	None	1/2
BRVS	k	Branch if Overflow Flag is Set	if (V = 1) then PC ← PC + k + 1	None	1/2
BRVC	k	Branch if Overflow Flag is Cleared	if (V = 0) then PC ← PC + k + 1	None	1/2

[2] Source: ATmega328P Data Sheet http://www.atmel.com/dyn/resources/prod_documents/8161S.pdf Chapter 31 Instruction Set Summary

Mnemonics	Operands	Description	Operation	Flags	#Clocks
BRIE	k	Branch if Interrupt Enabled	if (I = 1) then PC ← PC + k + 1	None	1/2
BRID	k	Branch if Interrupt Disabled	if (I = 0) then PC ← PC + k + 1	None	1/2
BIT AND BIT-TEST INSTRUCTIONS					
SBI	P,b	Set Bit in I/O Register	I/O(P,b) ← 1	None	2
CBI	P,b	Clear Bit in I/O Register	I/O(P,b) ← 0	None	2
LSL	Rd	Logical Shift Left	Rd(n+1) ← Rd(n), Rd(0) ← 0	Z,C,N,V	1
LSR	Rd	Logical Shift Right	Rd(n) ← Rd(n+1), Rd(7) ← 0	Z,C,N,V	1
ROL	Rd	Rotate Left Through Carry	Rd(0)←C,Rd(n+1)← Rd(n),C←Rd(7)	Z,C,N,V	1
ROR	Rd	Rotate Right Through Carry	Rd(7)←C,Rd(n)← Rd(n+1),C←Rd(0)	Z,C,N,V	1
ASR	Rd	Arithmetic Shift Right	Rd(n) ← Rd(n+1), n=0..6	Z,C,N,V	1
SWAP	Rd	Swap Nibbles	Rd(3..0)←Rd(7..4),Rd(7..4)←Rd(3..0)	None	1
BSET	s	Flag Set	SREG(s) ← 1	SREG(s)	1
BCLR	s	Flag Clear	SREG(s) ← 0	SREG(s)	1
BST	Rr, b	Bit Store from Register to T	T ← Rr(b)	T	1
BLD	Rd, b	Bit load from T to Register	Rd(b) ← T	None	1
SEC		Set Carry	C ← 1	C	1
CLC		Clear Carry	C ← 0	C	1
SEN		Set Negative Flag	N ← 1	N	1
CLN		Clear Negative Flag	N ← 0	N	1
SEZ		Set Zero Flag	Z ← 1	Z	1
CLZ		Clear Zero Flag	Z ← 0	Z	1
SEI		Global Interrupt Enable	I ← 1	I	1
CLI		Global Interrupt Disable	I ← 0	I	1
SES		Set Signed Test Flag	S ← 1	S	1
CLS		Clear Signed Test Flag	S ← 0	S	1
SEV		Set Twos Complement Overflow.	V ← 1	V	1
CLV		Clear Twos Complement Overflow	V ← 0	V	1
SET		Set T in SREG	T ← 1	T	1
CLT		Clear T in SREG	T ← 0	T	1
SEH		Set Half Carry Flag in SREG	H ← 1	H	1
CLH		Clear Half Carry Flag in SREG	H ← 0	H	1
DATA TRANSFER INSTRUCTIONS					
MOV	Rd, Rr	Move Between Registers	Rd ← Rr	None	1
MOVW	Rd, Rr	Copy Register Word	Rd+1:Rd ← Rr+1:Rr	None	1
LDI	Rd, K	Load Immediate	Rd ← K	None	1
LD	Rd, X	Load Indirect	Rd ← (X)	None	2
LD	Rd, X+	Load Indirect and Post-Inc.	Rd ← (X), X ← X + 1	None	2
LD	Rd, - X	Load Indirect and Pre-Dec.	X ← X - 1, Rd ← (X)	None	2
LD	Rd, Y	Load Indirect	Rd ← (Y)	None	2
LD	Rd, Y+	Load Indirect and Post-Inc.	Rd ← (Y), Y ← Y + 1	None	2
LD	Rd, - Y	Load Indirect and Pre-Dec.	Y ← Y - 1, Rd ← (Y)	None	2
LDD	Rd,Y+q	Load Indirect with Displacement	Rd ← (Y + q)	None	2
LD	Rd, Z	Load Indirect	Rd ← (Z)	None	2
LD	Rd, Z+	Load Indirect and Post-Inc.	Rd ← (Z), Z ← Z+1	None	2
LD	Rd, -Z	Load Indirect and Pre-Dec.	Z ← Z - 1, Rd ← (Z)	None	2
LDD	Rd, Z+q	Load Indirect with Displacement	Rd ← (Z + q)	None	2
LDS	Rd, k	Load Direct from SRAM	Rd ← (k)	None	2
ST	X, Rr	Store Indirect	(X) ← Rr	None	2
ST	X+, Rr	Store Indirect and Post-Inc.	(X) ← Rr, X ← X + 1	None	2
ST	- X, Rr	Store Indirect and Pre-Dec.	X ← X - 1, (X) ← Rr	None	2
ST	Y, Rr	Store Indirect	(Y) ← Rr	None	2
ST	Y+, Rr	Store Indirect and Post-Inc.	(Y) ← Rr, Y ← Y + 1	None	2
ST	- Y, Rr	Store Indirect and Pre-Dec.	Y ← Y - 1, (Y) ← Rr	None	2
STD	Y+q,Rr	Store Indirect with Displacement	(Y + q) ← Rr	None	2
ST	Z, Rr	Store Indirect	(Z) ← Rr	None	2
ST	Z+, Rr	Store Indirect and Post-Inc.	(Z) ← Rr, Z ← Z + 1	None	2
ST	-Z, Rr	Store Indirect and Pre-Dec.	Z ← Z - 1, (Z) ← Rr	None	2
STD	Z+q,Rr	Store Indirect with Displacement	(Z + q) ← Rr	None	2
STS	k, Rr	Store Direct to SRAM	(k) ← Rr	None	2
LPM		Load Program Memory	R0 ← (Z)	None	3
LPM	Rd, Z	Load Program Memory	Rd ← (Z)	None	3
LPM	Rd, Z+	Load Program Memory and Post-Inc	Rd ← (Z), Z ← Z+1	None	3
SPM		Store Program Memory	(Z) ← R1:R0	None	-
IN	Rd, P	In Port	Rd ← P	None	1
OUT	P, Rr	Out Port	P ← Rr	None	1
PUSH	Rr	Push Register on Stack	STACK ← Rr	None	2

Mnemonics	Operands	Description	Operation	Flags	#Clocks
POP	Rd	Pop Register from Stack	Rd ← STACK	None	2
MCU CONTROL INSTRUCTIONS					
NOP		No Operation		None	1
SLEEP		Sleep	(see specific descr. for Sleep function)	None	1
WDR		Watchdog Reset	(see specific descr. for WDR/timer)	None	1
BREAK		Break	For On-chip Debug Only	None	N/A

Note: 1. These instructions are only available in ATmega168PA and ATmega328P.